A NECESSARY AND SUFFICIENT CONDITION FOR THE UNIQUENESS OF MINIMUM SPANNING TREE

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SARI

Dengan menggunakan relasi samar sebagai sudut pandang, dalam tulisan ini dikembangkan suatu sifat fundamental dari penutup transitif min-maks suatu disimilaritas, dalam hubungannya dengan ultrametrik sub-dominan. Sifat tersebut memungkinkan kita merumuskan dan membuktikan suatu syarat cukup dan perlu agar suatu disimilaritas memiliki pohon kerangka minimum yang tunggal. Apabila tidak tunggal, sifat itu dapat menjadi landasan untuk menentukan semua pohon kerangka minimum lokal.

ABSTRACT

We develop a fundamental property of min-max transitive closure of a dissimilarity, considered as a fuzzy relation, in connection with its subdominant ultrametric. This will enable us firstly to derive a necessary and sufficient condition for the uniqueness of its minimum spanning tree, and secondly to find all possible local minima.

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The concept of minimum spanning tree (MST) was originally developed in the field of graph theory. In the last two decades we see its widespread use in many disciplines such as biology, social science, economy, antrophometry and general taxonomy [1], data analysis [2], regression analysis [4] and [5], computer science [6], networking [8], and multivariate and clustering analysis [10].

The ability to detect the uniqueness of MST is the first great problem for statisticians in using MST for their statistical analysis. The complexity of statistical analysis of a dissimilarity data matrix depends on the uniqueness of its MST (see [5], and [10]). Unfortunately, as far as we know (see [1], [3], [7], [8], [9], and [11]), there is no algorithm that can detect the uniqueness of MST. The second great problem is the fact that only one MST can be given by all existing algorithms, even for the case where there are actually more than one MST. This fact can also be found in theoritical literatures (see [6], [9], and [11]). In these two circumstances, Proposition 1, Proposition 2 and Proposition 3 are the main result of this work. In particular, Proposition 3 provides us with two fundamental results:

1. A necessary and sufficient condition for the uniqueness of MST.
2. An algorithm for constructing the union of all possible MSTs which, if it is not unique, gives us all MSTs.

The basic problem in this paper, firstly is to propose and to show a necessary and sufficient condition for uniqueness of MST of a dissimilarity. Secondly is to give an algorithm for finding all possible MSTs, if there are actually more than one MST. For this purpose we consider a dissimilarity as a fuzzy relation. Some basic concepts can be found in [3].

Suppose R a fuzzy relation on a set I; Card(I) = n and for every (x, y) in I x I we have 0 ≤ μR(x, y) < ∞, where μR is a membership function on R. In this paper we develop a fundamental property of min-max transitive closure R' of R, where R is a dissimilarity, in connection with subdominant ultrametric. Another representation of R' can be seen in [7] and [11]. The relationship between sub-dominant ultrametric and minimum spanning tree such as shown in [2], [6], and [9] will be exploited in order to derive a necessary and sufficient condition for the uniqueness of minimum spanning tree in a dissimilarity.

2. TRANSITIVE CLOSURE

A fuzzy relation R on I is called max-min transitive if for all x, y, z in I we have

μR(x, z) ≤ V_y {μR(x, y) ∧ μR(y, z)}

It is known that max-min transitivity of R can be verified through max-min composition o. If R'^2 = R o R is a fuzzy relation defined by
Theorem 1  \[ R^{02} = R, \text{ then } R \text{ is max-min transitive.} \]

Theorem 2  \[ R \text{ is max-min transitive if and only if } R^{02} \subseteq R. \]

Since \( I \) is finite, the max-min transitive closure \( R^- \) of \( R \) has the following representation:

\[ R^- = R \cup R^{02} \cup R^{03} \cup \ldots \cup R^k, \]

for an integer \( k; 1 \leq k \leq n \), where \( R^k = R \circ R \circ \ldots \circ R \), \( k \) times max-min composition \( \circ \) of \( R \).

Like \( R^- \), the min-max transitive closure \( R^+ \) can be written as

\[ R^+ = R \cap R^{*2} \cap R^{*3} \cap \ldots \cap R^k \]

for an integer \( k, 1 \leq k \leq n \), where \( R^k = R \ast R \ast \ldots \ast R \), \( k \) times min-max composition \( \ast \) of \( R \)

\[ \mu_{R^*} (x, z) = \bigwedge_y \{ \mu_R(x, y) \lor \mu_R(y, z) \} \]

If \( R^- \) and \( R^+ \) represent respectively the complement of \( R \) and min-max transitive closure of \( R \), it can be shown that \( R^- = R^+ \) and \( R^+ = R^+ \). Hence by De Morgan's rule we have \( R^+ = R^{cc} \) and \( R^- = R^{cc} \). These equalities enable us to work with either \( R^- \) or \( R^+ \). Although those representations are very useful, but it is still not comfortable to work with. The following alternative form which is more convenient for constructing \( R^- \) is given in [7].

Theorem 3  Suppose that \( R \) is a fuzzy relation on \( I \). Let

\[ l^*(x, y) = \bigvee_{c \in \zeta} l(c), \text{ where } \]

a. \( \zeta = \{ c \mid c = (x = x_{i_1}, x_{i_2}, \ldots, x_{i_r} = y) \} \) is a chain from \( x \) to \( y \)

b. \( l(c) = \mu_R((x_{i_1}, x_{i_2}) \cap \mu_R(x_{i_2}, x_{i_3}) \cap \ldots \cap \mu_R(x_{i_r-1}, x_{i_r}). \)

Then \( \mu_{R^*} (x, y) = l^*(x, y), \) for all \( x \) and \( y \) in \( I \).

In practice this theorem is still difficult to be implemented. In the next section we will restrict our discussion in the case where \( R \) is a dissimilarity and we derive a fundamental property in connection with its sub-dominant ultrametric in order to construct simple computation.
MIN-MAX TRANSITIVE CLOSURE AND SUB-DOMINANT ULTRAMETRIC

Suppose R a dissimilarity associated to dissimilarity index d on I. Hence R is a symmetric and anti-reflexive fuzzy relation, where \( \mu_R(x, y) = d(x, y) \) for all x and y in I. In the following proposition we show that, in this case, \( R^+ \) has a very convenient representation.

**Proposition 1** If R is a dissimilarity on I, then \( R^+ = R^k \) for an integer k; 1 \( \leq k \leq n \).

**Proof**

We know that for an integer \( k; 1 \leq k \leq n \),

\[
R^+ = R \cap R^+ \cap \ldots \cap R^k
\]

Now we show that the right hand side is equal to \( R^k \).

By definition,

\[
\mu_{R^+}(x, z) = \bigvee_{y} \{\mu_R(x, y) \lor \mu_R(y, z)\},
\]

for all x, y and z in I. Especially if y = z, then

\[
\mu_{R^+}(x, z) \leq \mu_R(x, z) \lor \mu_R(z, z)
\]

But \( \mu_R(z, z) = 0 \). Hence,

\[
\mu_{R^+}(x, z) \leq \mu_R(x, z)
\]

for all x and z in I or \( R^2 \subseteq R \). In general we have

\[
R^{2k} \subseteq \ldots \subseteq R^2 \subseteq R^k \subseteq R.
\]

It implies that \( R^+ = R^k \).

Now we show a fundamental property of \( R^+ \) in connection with sub-dominant ultrametric (SDU) of dissimilarity R.

**Proposition 2** If R is a dissimilarity on I, then \( R^+ \) is the SDU of R.

**Proof**

Theorem 3 tells us that

\[
\mu_{R^+}(x, y) = \mu_{R^2}(x, y) = \max_{C \in \mathcal{C}} I(C)
\]

where \( C = (x = x_1, x_2, \ldots, x_r = y) \) is a chain from x to y.

If \( \alpha = \mu_{R^+}(x, x) \) for all x in I, then
\[
\mu_{R^+}(x, y) = \max \{\min_k \{\mu_R(x_{i_k}, x_{i_{k+1}})\}\}
\]

\[
= \max \{\min \{\mu_R(x_{i_1}, x_{i_2}), \ldots, \mu_R(x_{i_{r-1}}, x_{i_r})\}\}\]

\[
= \max \{\mu_R(x_{i_1}, x_{i_2}), \ldots, \mu_R(x_{i_{r-1}}, x_{i_r})\}\]

\[
= \alpha - \min \{\alpha - \{\alpha - \max \{\mu_R(x_{i_1}, x_{i_2}), \ldots, \mu_R(x_{i_{r-1}}, x_{i_r})\}\}\}
\]

This equality implies that:

\[
F^*(x, y) = \min \{\mu_R(x_{i_1}, x_{i_2})\}\]

Now we will show that \(R^+\) is the SDU of \(R\).

i. It is clear that \(\mu_{R^+}(x, y) \leq \mu_R(x, y)\) for all \(x\) and \(y\) in \(I\), since \(R^+ = R^\ast \subseteq R\).

ii. If \(C = (x = x_{i_1}, x_{i_2}, \ldots, x_{i_r} = y)\) is a chain from \(x\) to \(y\), we note that \(L(C) = \max_k \{\mu_R(x_{i_k}, x_{i_{k+1}})\}\).

Suppose that \(C_1\) is a chain from \(x\) to \(y\) and \(C_2\) is a chain from \(y\) to \(z\), such that \(\mu_{R^+}(x,y) = L(C_1)\) and \(\mu_{R^+}(y,z) = L(C_2)\).

Suppose also that \(C_3\) is a chain from \(x\) to \(z\), constructed from \(C_1\) and \(C_2\) such that:

\[L(C_3) = \max \{L(C_1), L(C_2)\}\]

In this case,

\[L(C_3) = \max \{\mu_{R^+}(x,y), \mu_{R^+}(y,z)\}\]

and we have,

\[\mu_{R^+}(x,z) = \min_{C \in \Omega} L(C) \leq L(C_3)\]

\[\leq \max \{\{\mu_{R^+}(x,y), \mu_{R^+}(y,z)\}\}\]

It implies that \(R^+\) is an ultrametric on \(I\).

iii. Suppose that \(U\) is the USD of \(R\). Now we show that \(U = R^+\).

Consider a chain \(C_1 = (x = x_{i_1}, x_{i_2}, \ldots, x_{i_r} = y)\) from \(x\) to \(y\) where \(\mu_{R^+}(x, y) = L(C_1)\). Then,
a. \( \mu_U(x,y) \leq \max \{ \mu_U(x,z), \mu_U(y,z) \} \) for all \( x, y \) and \( z \) in \( I \), because \( U \) is an ultrametric. Especially,
\[
\mu_U(x,y) \leq \max \{ \mu_U(x, x_k), \mu_U(x_{i_k}, y) \}
\]
for all \( k = 1, 2, \ldots, r \). Hence,
\[
\mu_U(x,y) \leq \max \{ \mu_U(x, x_{i_2}), \mu_U(x_{i_2}, y) \}
\]
\[
\leq \max \{ \mu_U(x, x_{i_2}), \max \{ \mu_U(x_{i_2}, x_{i_3}), \mu_U(x_{i_3}, y) \} \}
\]
\[
\leq \max \{ \mu_U(x, x_{i_2}), \mu_U(x_{i_2}, x_{i_3}), \mu_U(x_{i_3}, y) \}
\]
In general we have
\[
\mu_U(x,y) \leq \max \{ \mu_U(x, x_{i_k}), \mu_U(x_{i_k}, y) \}, \ 1 \leq k \leq r - 1.
\]

b. \( U \) is the USD of \( R \). Then by definition, \( U \subseteq R \) or
\[
\mu_U(x,y) \leq \mu_R(x,y), \ \text{for all } x \text{ and } y \text{ in } I.
\]
From a and b, we have;
\[
\mu_U(x,y) \leq \max \{ \mu_R(x, x_{i_k}), \mu_R(x_{i_k}, y) \}, \ 1 \leq k \leq r - 1.
\]
\[
\leq L(C_1) = \mu_{R^+}(x,y).
\]
or
\[
\mu_U(x,y) \leq \mu_{R^+}(x,y).
\]
It has been shown that \( R^+ \) is an ultrametric and \( U \) is the SDU of \( R \). Hence the inequality \( \mu_U(x,y) \leq \mu_{R^+}(x,y) \) gives us \( \mu_U(x,y) = \mu_{R^+}(x,y) \) or \( U = R^+ \).

4. SUB-DOMINANT ULTRAMETRIC AND MINIMUM SPANNING TREE

Through the notion of sub-dominant ultrametric, in this section we will show a necessary and sufficient condition for the uniqueness of minimum spanning tree. Suppose \( M \) is a minimum spanning tree of dissimilarity \( R \) defined by a dissimilarity index \( d \). If \( i \) and \( j \) are arbitrary two vertices in \( M \), and \( (i = x_1, x_2, \ldots, x_r = j) \) is the chain from \( i \) to \( j \) in \( M \), we know that the distance \( d \) between \( i \) and \( j \) given by
\[
\delta(i,j) = \max_k d(x_k, x_{k+1})
\]
is the SDU of \( R \). Hence
\[ \mu_{R^*}(i, j) = \delta(i, j) \]
\[ = d(x_{k_0}, x_{k_0+1}) \]

for a positive integer \( k_0 \). This equality and the above propositions show that the number of zero entries of \((R - R^*)\), subtraction of two matrices in the usual sense, determines the uniqueness of its minimum spanning tree. More specifically we have the following proposition.

**Proposition 3** Dissimilarity \( R \) has a unique minimum spanning tree if and only if the number of zero entries in the lower (or upper) triangle matrix of \((R - R^*)\) below (or above) diagonal, is equal to \((n-1)\).

If in a dissimilarity there are more than one MSTs, then we can find all MSTs by inspecting zero entries of lower (or upper) triangle matrix of \((R - R^*)\) below (or above) diagonal; we delete all unnecessary zero entries.

5. CONCLUDING REMARK

The ability to detect the uniqueness of MST and the fact that only one MST can be given by all existing algorithms, even for the case where there are actually more than one MST, is the great problem for statisticians in using MST for their statistical analysis. We have handled this problem through the notion of fuzzy relation. There are three propositions resulted in this work; Proposition 1, Proposition 2 and Proposition 3. In particular, Proposition 3 provides us with two fundamental results:

1. A necessary and sufficient condition for the uniqueness of MST.
2. An algorithm for constructing the union of all possible MSTs which, if it is not unique, gives us all MSTs.

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7. REFERENCES


