

Dynamic Overlapping Clustering for Wireless Sensor Networks Based-on Particle Swarm Optimization

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Abstract. In the recent years, Wireless Sensor Network (WSN) has been one of the most interesting research topics because of its flexibility and many potential applications. However, in the applications, there are still resources constraints, including: energy, computation, and bandwidth. It is believed that clustering is the best solution for the need of energy efficiency and scalability. In order to reach the high level of energy efficiencies, mostly, the clustering algorithms avoid the possibility of overlap between clusters. But in fact, there are several applications that need the occurrence of overlaps between clusters. In this paper, we propose a Particle Swarm Optimization (PSO)-based Clustering algorithm that has capability to control the overlap between clusters but still it has an ability to reach energy efficiency. PSO is chosen because it has a light computation and can quickly reach convergence. This proposed algorithm performance is analytically and experimentally compared with clustering on LEACH. The result of the test shows that this proposed algorithm has a capability to control the rate of overlapping degree linearly. The testing on the PSO for clustering also shows the better performance than on LEACH, although there are a few problems related to its complexity.

Keywords: clustering; dynamic; energy-efficient; overlapping; PSO; WSN.

1 Introduction

In the last decade, Wireless Sensor Networks (WSN) has been attracting many researchers. It has been recognized a new means for various needs of application monitoring and data collections. The major advantages of WSN are its small dimension and that it doesn't require wiring, either for data transmissions or powers [1]. Both factors make it flexible to be applied in any kind of environments.

WSN can autonomously operate without any human interventions directly. Thus, it is very appropriate to be applied in dangerous environments. Sensor is held fairly and randomly in sensing-area, and then followed by a process of auto-configuration and auto-operation to take, to process, and then to send the

data to the monitoring center [1]. In many applications, the large and dense population is needed to cover the extent area and increase the accuracy of the data retrieval. Scalability and energy-aware strategies are indispensable for this type of environment.

Many researchers have shown that a mechanism of grouping nodes, called clustering, which is performed dynamically, has a capacity to solve the scalability and energy efficiency problems in WSN. Clustering is an algorithm that command network to create cluster formation autonomously by choosing nodes that has role as a cluster-head (CH) and determine the member nodes that are affiliated to each CH. Clustering algorithm is run periodically to maintain the load balance among the nodes. It means that cluster formation is changed periodically so that the role as a CH is rotated alternately among all nodes.

Clustering algorithms has been proposed by many researchers. LEACH [2] is one of clustering algorithms for WSN that was proposed in earlier time and it has been a comparison for the next proposals. The goal of this algorithm is to increase the network lifetime by using signal strength as parameter in determining the formation. The nature of LEACH algorithm is "distributed" in which the algorithm is executed autonomously at each node. The LEACH's proposal is followed by numbers of the next algorithm proposals, in which, each of them has its advantages and disadvantages. The more comprehensive review about clustering algorithm for WSN is described in paper [3].

For the sake of efficiency, in the most proposals of clustering algorithms, each node belongs only to one cluster. However, in several purposes and applications, several nodes in the network are required to affiliate to more than one cluster so the overlapping among clusters occurs. Several applications that require overlapping clusters are inter-cluster routing [4], node localization [5], and time synchronization protocols [6]. The overlapping is also required in order to obtain the modal parameter estimation more accurately in the system of Structural Health monitoring (SHM) [7].

Despite there have been many proposals for clustering algorithm, as far as we know, there is only one algorithm that specifically controls the overlapping between clusters, which is KOCA [8]. KOCA is a clustering algorithm that has capability to control overlapping clustering at certain level. Despite it is proved that KOCA has a good performance, it is still static. It means that the cluster formation doesn't change overtime. The roles as CH are only played by certain nodes so that the nodes are receiving more load than others. Therefore it will more quickly run out of energy. To overcome this, it needs an algorithm that is able to control the overlapping levels of energy consumption and also to maintain the balance of energy consumptions among nodes.

Moreover, most clustering algorithms are fully distributed, in which decisions are made independently by each node based on the knowledge of the condition itself and the condition of the surrounding nodes. In terms of data processing load, the distributed strategy will result in a very lightweight computation, thus it would be very suitable for WSN. However, the major disadvantage of distributed algorithm is that it is difficult to achieve optimal conditions because it uses limited information and doesn't make use of the optimization algorithms that are commonly used.

Several researchers has tried to apply some optimization techniques in clustering algorithms, such as LEACH-C, the centralized version of LEACH who uses simulated annealing (SA) optimization [2], and Tabu Search based Clustering [9],[10]. The uses of SA, Tabu Search, or other Heuristic Optimization in solving the clustering problem are based on the fact that clustering problem is an NP-hard problem that is very difficult to resolve mathematically. Optimization-based clustering is able to produce better performance than non-optimized. However, the optimization-based algorithm is a high computational process that requires amount of large energy resources and processing. Commonly, the optimization-based clustering algorithm is centralized because the algorithm is placed in the center, base-station (BS), which is not limited in resources.

Centralized strategy requires the transmission of information about the current state of each node to BS. It means that there is some additional information must be transmitted beyond the sensing data. If the BS is far from the network, another effect that emerges was the delay due to the transmission. To overcome this problem, it is crucial to select the lightweight optimization method to run at the node level of WSN.

Among many optimization methods, it is believed that the Particle Swarm Optimization is very suitable for WSN because it is easily implemented, efficiently computed, and can quickly achieve convergent [11]. PSO has been proposed by researchers as the solution of the problems in WSN, such as: deployment, localization, and data aggregation. PSO is also recognized for its use in clustering [12],[13] and it is showing good performance. However, to more closely realized, the use of PSO for the clustering needs to be investigated more deeply, including about the influence of PSO parameters on the clustering performances, when it is on the time of achieving the convergence, on the processing time, etc. that there are not discussed in both papers.

The work presented in this paper has three contributions:

1. We propose the energy efficient PSO-based Clustering algorithm with the ability to control the level of overlap between clusters.

Overlapping control mechanism that we propose is based on the threshold value of the distances comparison between the nodes to the main CH and the nodes to the secondary CH(s). The performance of O-PSO clustering algorithm is also shown in this paper using simulation and compared with O-LEACH [14].

2. We analyze the influence of PSO parameter in solving the problem of clustering.

The parameters of PSO investigated here include: the influence of particles and iterations number on the solution accuracy, the influence of particles number of the initial data on the speed of convergence, the effect of iterations and the particles number on the energy efficiency. In this paper, we will also show the evaluations of the cluster formations produced, including the distributions of cluster members and cluster width.

3. We analyze mathematically the influence of the clusters uniformity on the energy consumption.

This analysis will correlate with the evaluation result PSO parameters as stated in the contribution number 2.

Furthermore, this paper is systematically arranged as follows: Chapter 2 explains about the formulation of clustering problem and the solution using PSO; Chapter 3 explains in detail about algorithm O-PSO clustering as we proposed. Chapter 4 shows the simulation results and the analysis. At last, Chapter 5 resumes our conclusion about of the work.

2 Clustering Problem on WSN and Its Solution Using PSO

The basic principle of clustering algorithm is the grouping of nodes in a network into several clusters, in which each cluster led by a CH. Using clustering algorithm, the network will autonomously select several nodes that act as CH and the member nodes affiliated to each CH. Unlike the commonly wireless networks, the goal of clustering in WSN is to increase energy more efficiently than network throughput. The optimization problem of clustering in WSN is to find the cluster formation in such a way as to produce very minimum energy consumption.

2.1 Energy Model

Commonly used in many papers, for example [2], energy consumption model in WSN used here has three parts, they are:

1. For sending an *l*-bit message with distance *d*, a node require energy

$$E_{Tx}(l,d) = \begin{cases} l. E_{elec} + l. \epsilon_{fS} d^2 f or d < d_0 \\ l. E_{elec} + l. \epsilon_{mp} d^4 f or d \ge d_0 \end{cases}$$
 (1)

2. For receiving an *l*-bit message, a node require energy

$$E_{Rx}(l) = l.E_{elec} \tag{2}$$

3. Energy for data aggregation is E_{DA}

Based on this model, Yan Jin et.al. [15] proposed a lemma about when a clustered network consume a fewer energy than a network without clustering. They claimed that clustering will be more efficient in energy consumption than direct transmission when a formation filled with the following condition [15]:

$$D \ge \begin{cases} \sqrt{\frac{E_{elec}}{\epsilon_{fs}}} for R < d_0 \\ \sqrt{\frac{E_{elec}}{\epsilon_{fs}}} for R \ge d_0 \end{cases}$$
 (3)

where *D* is the distance from CH to BS and *R* is the cluster radius. Assume the data is ideally aggregated and the energy for the aggregation process is ignored.

2.2 Impact of Clustering Uniformity on Energy Efficiency of WSN

Using energy model in [2] and by using the same assumption with in [15], the following exposure contains analysis that we are proposing in order to prove that a network with more uniform cluster formation will be more energy efficient than with the less uniform. The uniformity that is meant here is the similarity in the number of the cluster members in one formation. We conduct this analysis to prove our hypothesis that within the fairly member number, the load of each CH will considerably similar. Thus, there is no CH which gets the heavier load than other, which can result on the energy left-out more quickly.

Assuming that a cluster k contains N_k nodes with radius R_k and the distance from CH to BS is D_k , based on [2] and [15], we can determine that the energy consumption of cluster k will be

$$E_k = (2N_k - 1)lE_{elec} + l\epsilon_{mp}D_k^4 + (N_k - 1)l\epsilon_{fs}\frac{R_k^2}{2}$$
(4)

If node is deployed in the sensing area uniformly, the number of node in an area could be analogous to the wide of the area, therefore

$$N_k \approx R_k^2 \tag{5}$$

The influence of cluster uniformity on energy consumption can be analyzed by comparing two different networks in the cluster uniformity, as illustrated at Figure 1.

Total energy consumed by formation A is

$$E_A = E_{A1} + E_{A2} + E_{A3} \tag{6}$$

Based on Eq. (4), therefore

$$E_A = 3(2N - 1)lE_{elec} + \sum_{k=1}^{3} l\epsilon_{mp} D_k^4 + 3(N - 1)l\epsilon_{fs} \frac{R^2}{2}$$
 (7)

For formation B where $N_2 = 1.5N_1 = 1.5N$, and $N_3 = 0.5N_1 = 0.5N$ and with the reference to Eq. (5), we can calculate the radius square of the cluster B1, B2, and B3 respectively are R^2 , $1.5R^2$ and $0.5R^2$. Therefore, the total energy for formation B is

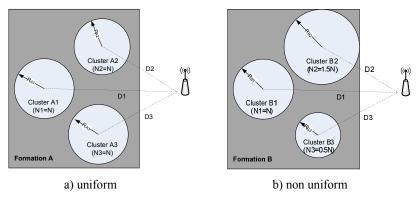


Figure 1 Two clusters formation with different uniformity.

$$E_{B} = E_{B1} + E_{B2} + E_{B3}$$

$$E_{B} = \{(2N - 1)(3N - 1)(N - 1)\} l E_{elec} +$$

$$\sum_{k=1}^{3} l \epsilon_{mp} D_{k}^{4} + \sum_{k=1}^{3} (N_{k} - 1) l \epsilon_{fs} \frac{R_{k}^{2}}{2}$$

$$E_{B} = (6N - 3) l E_{elec} + \sum_{k=1}^{3} l \epsilon_{mp} D_{k}^{4} +$$

$$(9)$$

$$(N-1)l\epsilon_{fs}\frac{R^2}{2} + (1.5N-1)l\epsilon_{fs}\frac{1.5R^2}{2} + (0.5N-1)l\epsilon_{fs}\frac{0.5R^2}{2}$$
(10)

$$E_B = 3(2N - 1)lE_{elec} + \sum_{k=1}^{3} l\epsilon_{mp}D_k^4 +$$

$${(N-1) + 1.5(1.5N-1)l\epsilon_{fs} + 0.5(0.5N-1)}l\epsilon_{fs} \frac{R^2}{2}$$
 (11)

$$E_B = 3(2N-1)lE_{elec} + \sum_{k=1}^{3} l\epsilon_{mp}D_k^4 +$$

$$\{N-1+2.25N-1.5+0.25N-0.5\}l\epsilon_{fs}\frac{R^2}{2}$$
 (12)

$$E_B = 3(2N - 1)lE_{elec} + \sum_{k=1}^{3} l\epsilon_{mp} D_k^4 + (3.5N - 3)l\epsilon_{fs} \frac{R^2}{2}$$
 (13)

From Eqs. (7) and (13), we can calculate that the difference of energy consumption between formation B and formation A is

$$E_B - E_A = \left\{ 3(2N - 1)lE_{elec} + \sum_{k=1}^3 l\epsilon_{mp} D_k^4 + (3.5N - 3)l\epsilon_{fs} \frac{R^2}{2} \right\} - \left\{ 3(2N - 1)lE_{elec} + \sum_{k=1}^3 l\epsilon_{mp} D_k^4 + 3(N - 1)l\epsilon_{fs} \frac{R^2}{2} \right\}$$
(14)

$$E_B - E_A = \{(3.5N - 3) - 3(N - 1)\}l\epsilon_{fs} \frac{R^2}{2} = 0.5Nl\epsilon_{fs} \frac{R^2}{2}$$
 (15)

$$E_B - E_A = 0.5Nl\epsilon_{fS} \frac{R^2}{2} \tag{16}$$

With N and R are always positive, then E_B in Eq. (15) will always be bigger than E_A , thus we can conclude that, with cluster that are not uniform, formation B will always consume more energy than formation A does with uniform cluster.

2.3 Optimization Problem Formulation for WSN Clustering

If the number of clusters is determined at the beginning to make N_c fixed, Eqs. (3) and (4) would generate a condition in which the amount of energy consumption is only determined by the distance of variable d. Both equations also shows that if $d_{toBS} \gg d_{toCH}$, then $E_{CH} \gg E_{nonCH}$. This condition shows that the nodes' role rotation as CH is very crucial to avoid imbalance in energy consumptions among nodes.

Based on this condition, the formulation of clustering optimization problem can be defined as follows [12]

Determine Kcluster-heads (CHs)and the related member nodes, to minimize:

$$f = \beta \times f_1 + (1 - \beta) \times f_2 \tag{17}$$

with f_1 is the maximum average Euclidean distance of nodes to their associated cluster heads and f_2 is ratio of total initial energy of all nodes in the network with the total energy of the CH candidates in current round. Both sub-functions are proportionally combined with weighting factor β . Function f_1 can be calculated using following formula:

$$f_1 = \max_{k=1,2,\dots,K} \left\{ \sum_{\forall n_i \in C_k} \frac{d(n_i, CH_k)}{|C_k|} \right\}$$
 (18)

where K is predetermined cluster number, C_k is the member of kth cluster, CH_k is the cluster-head of kth cluster, and d(x,y) is Euclidean distance between x and y. The second objective function can be formulated as follows:

$$f_2 = \frac{\sum_{i=1}^{N} E(n_i)}{\sum_{i=1}^{K} E(CH_k)} \tag{19}$$

where $E(n_i)$ and $E(CH_k)$ is residual energy of node i and cluster-head k respectively.

The problem uses the same assumptions those used in [2], that is: every node has the same probability to be a CH, every node can reach all other nodes and the BS, there is no communication interference between nodes and it is held randomly in 2D field, in which the location coordinate of each node is identified.

2.4 Finding Solution for Clustering Problem Using PSO

To address the problems in sub-chapter 2.3, we can use heuristic algorithms. One of them is Particle Swarm Optimization (PSO). Initially proposed by James Kennedy and Russell C. Eberhart [16], PSO is a population-based search algorithm, inspired by a social behavior of birds group.

There are some basic principles of PSO: firstly, PSO chooses S sets of solution candidates called as *particles*. Then, iteratively, each particle is moved to a new position with any particular *velocity*. Consider x_i is a position vector for particle i and v_i is a velocity for particle i. Each particle will be evaluated with an objective function f(x). In the case of minimization problem, we state that a particle is closer to the global solution if it can produce f(x) *smaller than* the other particles do. Velocity and position at iteration k is updated using Eqs. (20) and (21) and repeated until they reach a certain g_{best} value or have reached a number of iteration k_{max} that had been predetermined.

$$v_i(k) = w. v_i(k-1) + c_1 r_1(k) (p_{best_i} - x_i) + c_2 r_2(k) (g_{best_i} - x_i)$$
(20)

and

$$x_i(k) = x_i(k-1) + v_i(k)$$
 (21)

with w is inertia weight, c_1 and c_2 is learning factor constants, and r_1 and r_2 is random numbers uniformly distributed in [0,1].

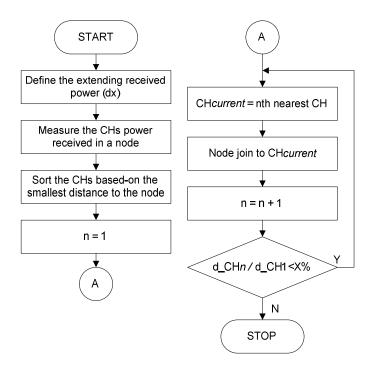


Figure 2 Flow chart of overlapping control mechanism.

In clustering problem, a particle consists of a set of node coordinate, nominated as CH candidate. Each particle is evaluated using objective function (7) until (9) for any iteration. If the formation of a particle in any iteration produces an f value bigger than the previous iterations values, therefore this formation will be saved as p_{best} . In any iteration, p_{best} of all particles are compared with each other to choose which particle producing the lowest f value, to be selected as the g_{best} . Variable of p_{best} and g_{best} is used to update the position and velocity of particles, as Eqs. (20) and (21). The PSO-based clustering is explained in detail in [12]

```
I. Set-up Phase
1. Compute average energy level of all nodes
2. CH candidates ←the nodes with energy level above the average
3. Initialize S particles to contain K randomly selected CH among the CH candidates
Repeat
      For each particle s_i, j = 1, 2, ..., S:
5.
           For each node n_i, i = 1,2,...,N:
              Calculate distance d(n_i, CH_k) between node n_i and all CH_k
6.
              Assign node n_i to CH_k where:
7.
                     d(n_i, CH_k) = \min_{\forall k=1,2,\dots,K} \{d(n_i, CH_k)\}
8.
              Assign node n_i to CH_vif:
                     d(n_i, CH_v) \le X\% of d(n_i, CH_k) v \ne k
9.
           End for
10.
          Calculate the fitness using equations (1)-(3)
          Find the personal best and global best
11.
12.
          Perform the position and velocity update
13.
          Limit the change in the particle's position value
14.
          Map the new updated position with closest (x,y) coordinates
15. End for
Until maximum number of iterations is reached
      Assign node n_i to CH_k where:
                                       d(n_i, CH_k) = \min_{\forall k=1,2,\dots,K} \{d(n_i, CH_k)\}
      Assign node n_i to other(s) CH_k if:
                             d(n_i, CH_k) \le X\% \text{ of } \min_{\forall k=1,2,\dots,K} \{d(n_i, CH_k)\}
II. Steady State Phase
Steady state operation for t=T_{round} second
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Figure 3 Pseudo code of O-PSO clustering.

3 Proposed Overlapping Clustering Algorithm Using PSO

In an energy efficient clustering algorithm, a non-CH node will choose to join a CH which has the closest position to it's after a set of CH has been decided. Provided that it is with the assumption that the channel characteristics between nodes are same so that the amount of energy used for data transmission is only determined by the distance. The same are also conducted in our proposed algorithm. A node will join to the closest CH candidate offered by a particle. A CH where a node first affiliated is called as primary CH.

Then, the node will join to other CH as a secondary CH as long as it is at a distance that is not farther than the threshold limit. A threshold value is calculated relatively to the node distance, in which the main CH is expressed in X%. The determination mechanism of primary and Secondary CH that produces overlapping between clusters is shown in detail at Figure 2. Then, the mechanism at Figure 2 is combined with the PSO-based clustering algorithm proposed in [7] to form a new algorithm which we refer as PSO-based Overlapping Clustering (O-PSO Clustering) Algorithm. Pseudo-code O-PSO Clustering is shown at Figure 3.

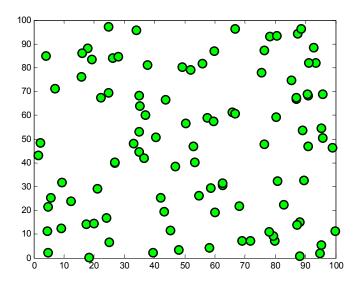
4 Simulation Results and Analysis

In this paper, the performance of clustering algorithm in energy efficiency is expressed as the number of the live nodes (N_{live}) in each *round*. We state that a node will be dead if the battery is empty (remaining energy = 0 Joule). This rule is usually used in WSN papers. The performance of energy efficiency of PSO-C is shown in the graphic and compared with the performances of LEACH and non-clustering. Meanwhile, the performance of O-PSO Clustering is evaluated for five variations of X value targets and compared with Overlapping LEACH (O-LEACH) [14].

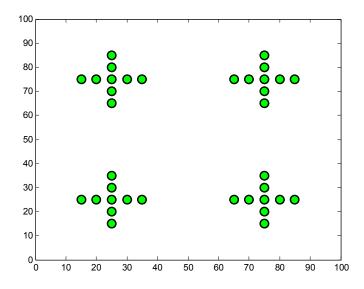
4.1 Experimental Setup

The use of LEACH and O-LEACH as the comparisons is based on the consideration that the assumptions used on O-PSO are same with the assumptions used on LEACH, it is beyond the fact that LEACH is very often used as comparison in clustering papers. Evaluations for both algorithms are performed by using Matlab simulation, in which we hold 100 nodes randomly in the area of 100m x 100m, as shown at Figure 4a.The life-time performance will be analyzed based on the uniformity the clusters on their number of members. The distribution of cluster is shown in the form of histogram for the first 30 rounds with the consideration that it is ensured there were no deadnodes in the round.

Especially for the testing on PSO optimization performance for clustering, we set the value of β =1, thus PSO just looks the cluster formation that produces minimal cluster diameter. Due to this requirement, we use deployment deterministic, as shown at Figure 3b, in which the cluster formation could be optimum if it is able to choose four CHs at the middle of each "+". The complete value of simulation parameter is shown in Table 1.



a) Random deployment (100 nodes uniformly deployed on 100m x 100m area)



b) Deterministic deployment(36 nodes deployed on 100m x 100m area)

Figure 4 Node deployment (BS not be displayed).

Parameter	Value
Network grid	From (0,0) to (100,100)
BS coordinate	(50,175)
Initial energy	2 J/battery
E_{elec}	50nJ/bit
\mathcal{E}_{fs}	10pJ/bit/m ²
$arepsilon_{mp}$	0.0013 pJ/bit/m ⁴
E_{DA}	5 nJ/bit/signal
Data packet size	500 bytes
Header packet size	25 bytes

Table 1 The Simulation Parameter.

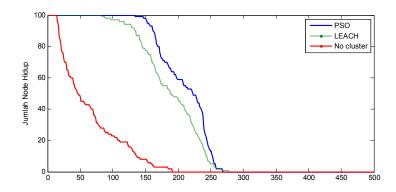


Figure 5 Lifetime performance of PSO-C compared to LEACH and no-clustering.

4.2 Evaluation of PSO on Solving Clustering Problem

The performance of PSO-C' energy efficiency compared with LEACH and non-clustering is shown at Figure 5. The figure shows two phenomena. Firstly, the clustered formation, either with PSO-C or LEACH, produces better performance than with the Non-clustering. Secondly, PSO-C has the better performance in energy consumption than LEACH. Both phenomena can be explained in following description.

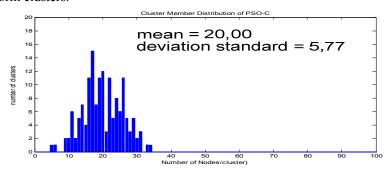
If the position of BS is far from CH, then clustering would consume less energy than non-clustering, as stated in Eq. (3). In this testing, we put BS at coordinate (50,175). By holding nodes in the field between (0,0) and (100,100), then the closest coordinate to BS is coordinate position of (50,100), that it is 75 in range from BS. Given the value of $E_{elec} = 50 \text{nJ/bit}$, $\varepsilon_{fs} = 10 \text{ pJ/bit/m}^2$, and $\varepsilon_{mp} = 0.0013 \text{ pJ/bit/m}^4$, thus:

$$\sqrt{\frac{E_{elec}}{\epsilon_{fs}}} = \sqrt{\frac{50^{-9}}{10^{-12}}} = 70.71 \le 75 \tag{22}$$

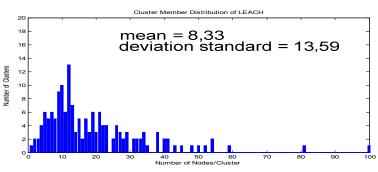
$$\sqrt[4]{\frac{E_{elec}/\epsilon_{mp}}{\epsilon_{mp}}} = \sqrt[4]{\frac{50^{-9}}{0.013.10^{-12}}} = 29.62 \le 75$$
 (23)

From the calculation (22) and (23), we can see that the distance CH-BS, D, on this deployment has fulfilled the minimal requirement in Eq. (3).

The second phenomenon that the PSO-C is more efficient in energy consumption than LEACH can be analyzed based on the uniformity property of their clusters. Figure 6 shows that the picture of cluster distribution on PSO-C is narrower than on LEACH. It means that PSO-C is more uniform than LEACH. This condition is suitable with sub-chapter 2.2, which states that a network with the more uniform clusters will consume less energy than the network with non-uniform clusters.



a) Cluster member distribution of PSO-Clustering



b) Cluster member distribution of LEACH

Figure 6 Distribution of member number on clusters.

The next evaluation is investigating the influence of particle number on the search errors, namely the difference between the solutions offered with the ideal solution. Figure 7 shows that each additional 10 particles would reduce the error estimation of CH location approximately 0.25 meters or cut about 10 iterations to reach convergence.

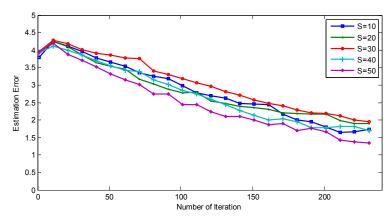


Figure 7 Impact of particle number to error estimation of CH position.

Another PSO parameter also evaluated is about the influence of particle initial quality. Figure 8 shows that the initial value is very influencing on the number of iterations required to achieve convergence especially for a small number of particles.

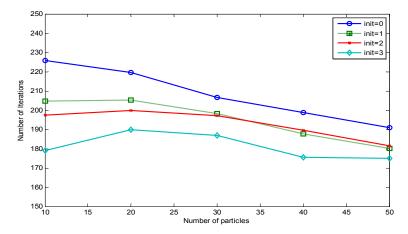


Figure 8 Impact of initial particle quality to convergence time.

Complexity is one of evaluations that are also important in testing an algorithm. The very simple way to measure the complexity of an algorithm is by measuring the time needed to run the algorithm. Figure 9 showed that the execution time of PSO-C is around a thousand times longer than LEACH. The figure showed that the number of iterations and particles will significantly increase the execution time.

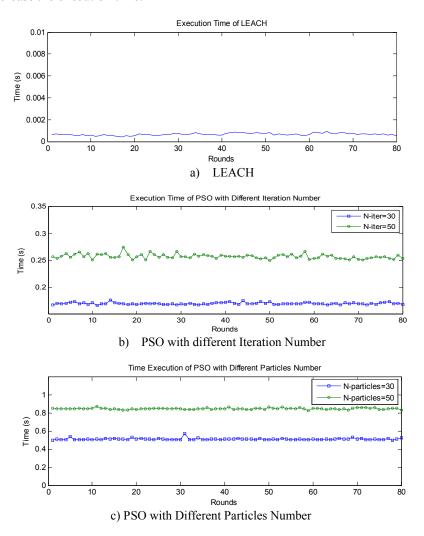


Figure 9 Execution time of the algorithm.

4.3 Evaluation of Overlapping PSO Clustering

The first performance indicator of O-PSO Clustering will be evaluated is the average of overlapping degree (OD). OD is the number of joint member nodes of two adjacent clusters. Figure 10 shows that the O-PSO clustering and O-LEACH produce an average OD value that rises linearly to the value. This phenomenon shows that the functions of OD control are working well for both algorithms. O-PSO has a greater slope than O-LEACH, which means that within the same X value, O-PSO is more efficient because it is able to achieve the greater OD average.

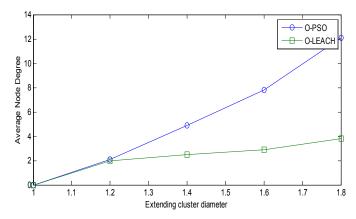


Figure 10 Overlapping degree (OD) of O-PSO compared to O-LEACH.

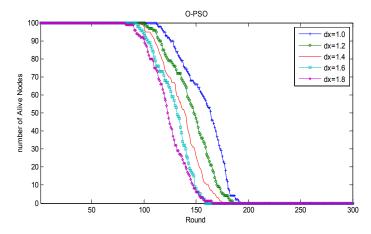


Figure 11 Network lifetime of O-PSO clustering.

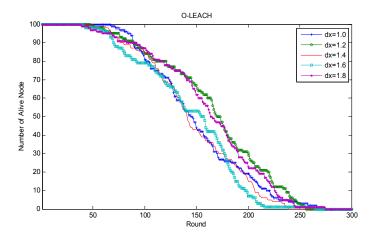


Figure 12 Network lifetime of O-LEACH.

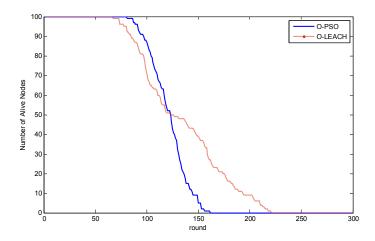


Figure 13 Network lifetime comparison: O-PSO vs O-LEACH.

The final evaluation is to test the influences of overlapping level on energy efficiency. The result of the testing on O-PSO and O-LEACH is shown at Figure 11 and Figure 12. From both figures, we can see that the time period spent for a node to die for the first time decreases less than 10 rounds for each 20% increase of X. It is understandable because several nodes is compelled to transmit messages to more than one CH. Figure 13 shows the comparison of performances between O-PSO and O-LEACH. The Figure shows that O-PSO is able to hold network's lifetime longer than O-PSO. It will happen continuously until the number of the live nodes reaching 50% of the original state.

5 Conclusion

This paper explains about a new proposal of PSO-based clustering for WSN that has a capability to control the level of overlapping between clusters, but still concerns on the energy efficiency level. This paper also shows the mathematical analysis about the influences of cluster uniformity on the level of energy consumption. The testing on this proposed algorithm shows the better performance than on O-LEACH. At the same parameter of threshold X, the O-PSO Clustering has capability to create the greater average of the overlapping level. O-PSO is also capable to keep the lower energy consumption. This paper also shows that PSO is very potential to be applied as a problem solving for clustering, despite there are still some problems to be solved, especially related to the algorithm complexities.

Acknowledgement

This research is carried out with the financial support of BPPS Scholarship from Indonesia Ministry of National Education. Some of the results of this paper have been presented in 2011 IEEE Region 10 Conference (TENCON 2011).

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